Most classical gates are one-way. After processing the input is not fully recoverable.
Truth-table for "and":

| $A$ | $B$ | $A$ and B |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

Only in case the output equals 1 we know that both inputs are 1 too. In case the output is 0 we cannot determine the state of inputs $A$ or $B$.

The NOT-gate is reversible:

| A | $\bar{A}$ |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |

Quantum gates need to be reversible. Each output must correspond to a specific input. The following are reversible quantum gates.

NOT


|  |  |
| :--- | :--- |
| I | 0 |
| 0 | 1 |
| 1 | 0 |

SWAP


Note: Swap only visible if $I_{1} \neq I_{2}$.
CNOT


The CNOT negates $I_{2}$ if $I_{1}=1$.

|  |  |  |  |
| :---: | :---: | :---: | :---: |
| $I_{1}$ | $I_{2}$ | $O_{1}$ | $O_{2}$ |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 |


| $I_{1}$ | $I_{2}$ | $O_{1}$ | $O_{2}$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 |

TOFFOLI


The Toffoli-gate negates $I_{3}$ if $I_{1}=I_{2}=1$.

| $I_{1}$ | $I_{2}$ | $I_{3}$ | $O_{1}$ | $O_{2}$ | $O_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 |

## FREDKIN



If $I_{1}=1$ the Fredkin-gate swaps $I_{2}$ and $I_{3}$. The effect is visible only if $I_{2} \neq I_{3}$.

| $I_{1}$ | $I_{2}$ | $I_{3}$ | $O_{1}$ | $O_{2}$ | $O_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 |

We can put this in other words.
The FREDKIN-gate swaps $I_{2}$ and $I_{3}$ if $I_{1}=1$, a controlled swap.
The TOFFOLI-gate negates $I_{3}$ if $I_{1}$ and $I_{2}$ both are 1 , a double controlled not.
The CNOT-gate negates $I_{2}$ if $I_{1}=1$, a controlled not.
The SWAP and the NOT do what they are expected to do.
Note: The gates are reversible because we have unique combinations on input-side and output-side.

## Example

For easier reading we name the lines $a, b, c, d$ and omit the distinctions between input and output.
We want a combination that

- $\quad$ swaps lines $b$ and $d$ if line $a=0$ and
- $\quad$ swaps lines $c$ and $d$ if line $a=1$ and
- leaves line $a$ untouched


| input |  |  |  | output |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| a | b | c | d | a | b | c | d |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 1 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

## Example

We want a combination that

- $\quad$ swaps lines $c$ and $d$ if line $a=b=1$


We need an auxiliary line $e$ to perform this.

| input |  |  |  | output |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| a | b | c | d | a | b | c | d |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

